

Bob



Alice



Pedersen Commitment

Prof Bill Buchanan, The Cyber Academy

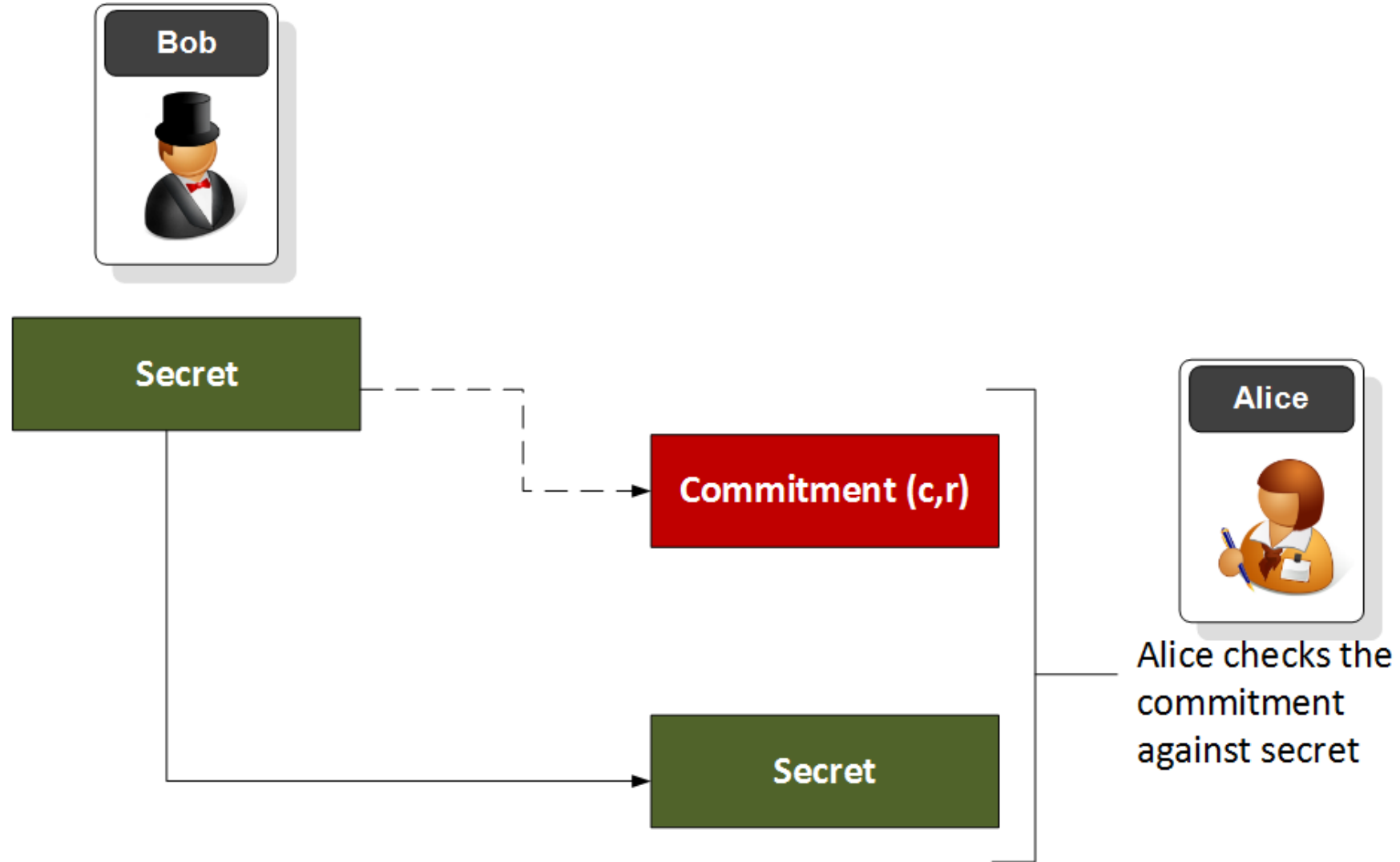
<http://asecuritysite.com>

Eve



**CYBER
ACADEMY**

Pedersen Commitment: Perfect Secrecy



Pedersen Commitment (Discrete Logs)

In the Pedersen commitment we take two large prime numbers (p and q) and we create a generator value (g) which is of the order of q and a subgroup of Z_p^* . Then s becomes a secret from 0 to Z_q , and we calculate:

$$h = g^s \pmod{p}$$

The sender now creates a commitment for a message (m) and with a random number (r):

$$c = g^m h^r \pmod{p}$$

The sender can then send c to the receiver. Next, the sender will then reveal m and r for the commitment, and the receiver verifies with:

$$c = g^m h^r \pmod{p}$$

If the values match, the receiver has proven the commitment.

We can also use as homomorphic encryption, and where we multiply commitments, and which is equivalent to adding the secret values.

Pedersen Commitment (Disc

In the Pedersen commitment we take two large values (g) which is of the order of q and a subgroup H of order p . We then calculate:

$$h = g^s \pmod{p}$$

The sender now creates a commitment for a message m and a random value r :

$$c = g^m h^r \pmod{p}$$

The sender can then send c to the receiver. The receiver can verify the commitment by checking if c is in the subgroup H and the receiver verifies with:

$$c = g^m h^r \pmod{p}$$

If the values match, the receiver has proven that the sender committed to the message m .

We can also use it as homomorphic encryption, where we can add commitments together to get a commitment to the sum of the secret values.

```
class verifier:
    def setup(self, security):
        p = number.getPrime(2 * security, Random.new().read)
        q = 2*p + 1

        g = number.getRandomRange(1, q-1)
        s = number.getRandomRange(1, q-1)
        print "Secret value:\t",s
        h = pow(g,s,q)

        param = (p,q,g,h)
        print "p=",p
        print "q=",q
        print "g=",g
        print "h=",h

        return param
```

Pedersen Commitment (Disc

In the Pedersen commitment we take two large values (g) which is of the order of q and a subgroup H of order r and calculate:

$$h = g^s \pmod{p}$$

The sender now creates a commitment for a message m and a decommitment r :

$$c = g^m h^r \pmod{p}$$

The sender can then send c to the receiver. The receiver can verify the commitment by checking if c is in the subgroup H and the receiver verifies with:

$$c = g^m h^r \pmod{p}$$

If the values match, the receiver has proven that the sender committed to the message m .

We can also use Pedersen commitment as homomorphic encryption, where the commitment of the sum of two messages is equal to the product of their individual commitments, i.e. adding the secret values.

```
class Verifier:
    security = 80
    msg1 = 1
    msg2 = 2

    def setup(self, security):
        p = if (len(sys.argv)>1):
            msg1=int(sys.argv[1])
        q = if (len(sys.argv)>2):
            msg2=int(sys.argv[2])

        g = v = verifier()
        p = prover()
        s =

    print param = v.setup(security)

    h = c1, r1 = p.commit(param, msg1)
        c2, r2 = p.commit(param, msg2)

    param = v.add(param, c1, c2)

    print "\nMsg1:",msg1
    print "Msg2:",msg2

    print "\nc1,r1:",c1,",",r1
    print "c2,r2:",c2,",",r2
    print "\nWe can now multiply c1 and c2, which is same as adding Msg1 and Msg2"
    print "\nCommitment of adding (Msg1+Msg2):\t",addCM

    result1 = v.open(param, c1, msg1, r1)
    result2 = v.open(param, c2, msg2, r2)

    print "\nResult of verifying c1:\t\t",result1
    print "Result of verifying c2:\t\t",result2

    result = v.open(param, addCM, msg1 + msg2, r1, r2)

    print "Result of verify Msg+Msg2:\t",result
```

Pedersen Commitment (Disc

```
class Verifier:
    security = 80
    msg1 = 1
    msg2 = 2
    def set(self, msg1, msg2):
        self.msg1 = msg1
        self.msg2 = msg2
    def verify(self, c1, c2):
        return (c1[0] == self.msg1 and c1[1] == self.msg2)
```

Secret value: 1600218503816068609616068262186085211385912913449

p= 835633126610936170404251455173303609923742171077

q= 1671266253221872340808502910346607219847484342155

g= 853879156962728674489578584402624767640115848839

h= 312816528258481903392461143384659987330939968504

Msg1: 70

Msg2: 70

c1,r1: 931432606187308599362760090383550998212817597419 , 1331540638789558839751757756218591899240229287857

c2,r2: 377312755894089489173415862148929013782547633796 , 1062296325886690334534931368033516581934275153224

Msg2"

We can now multiply c1 and c2, which is same as adding Msg1 and Msg2

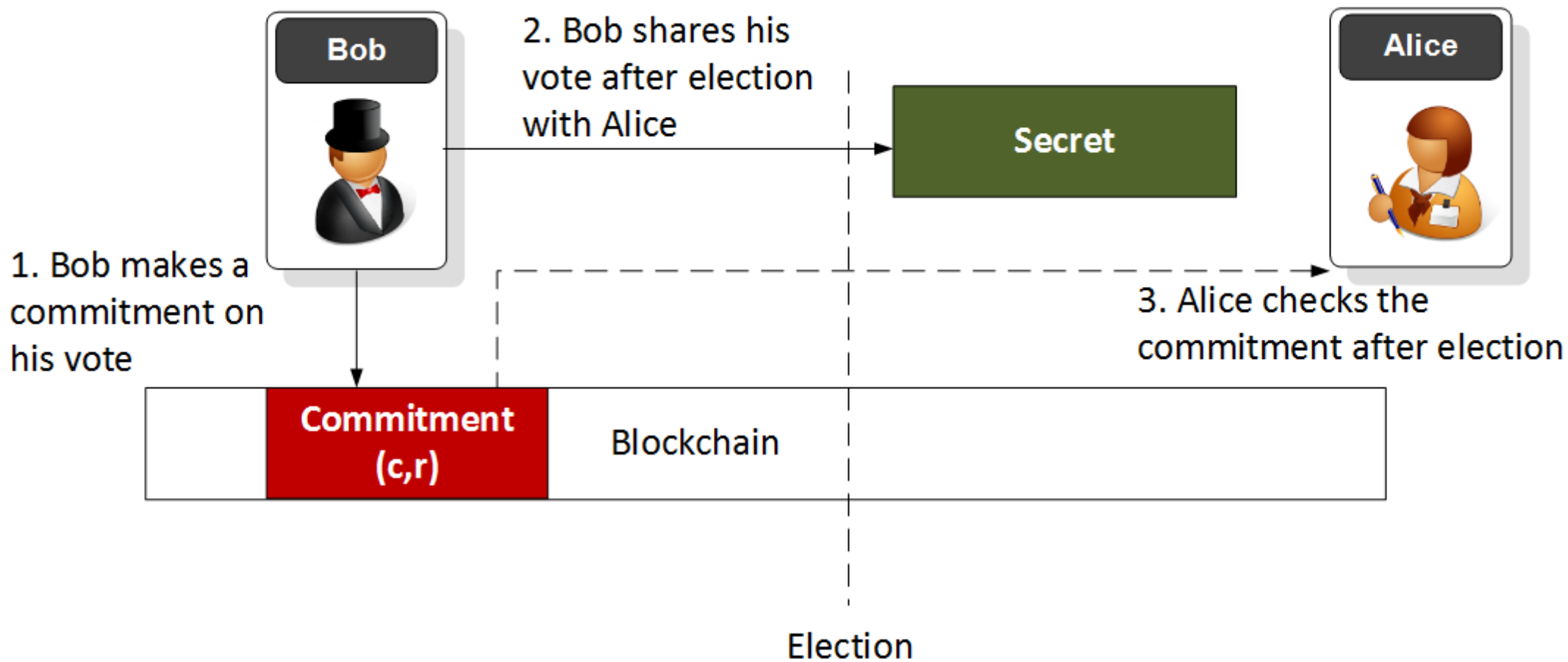
Commitment of adding (Msg1+Msg2): 943416350824651526788742276019896656762684796624

Result of verifying c1: True

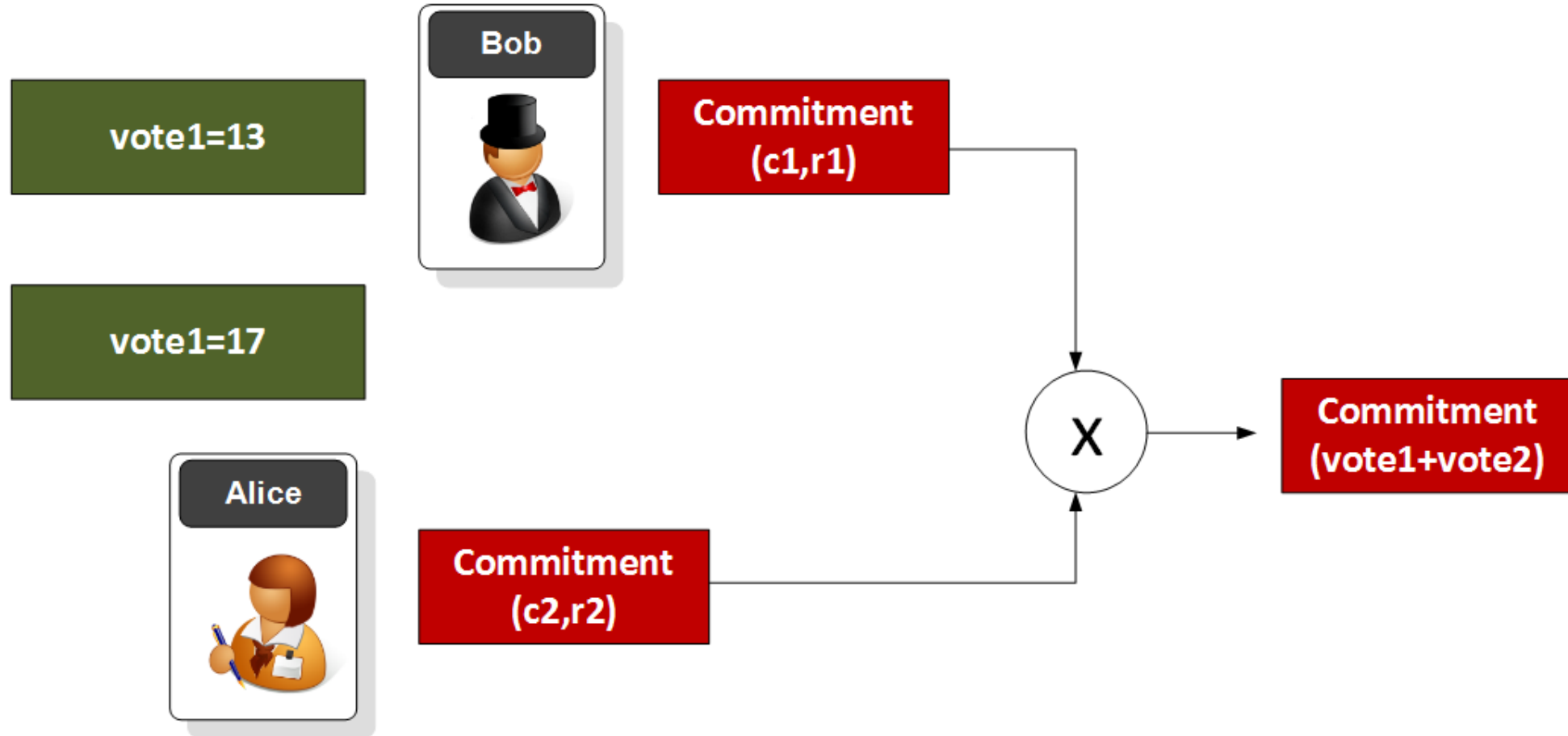
Result of verifying c2: True

Result of verify Msg+Msg2: True

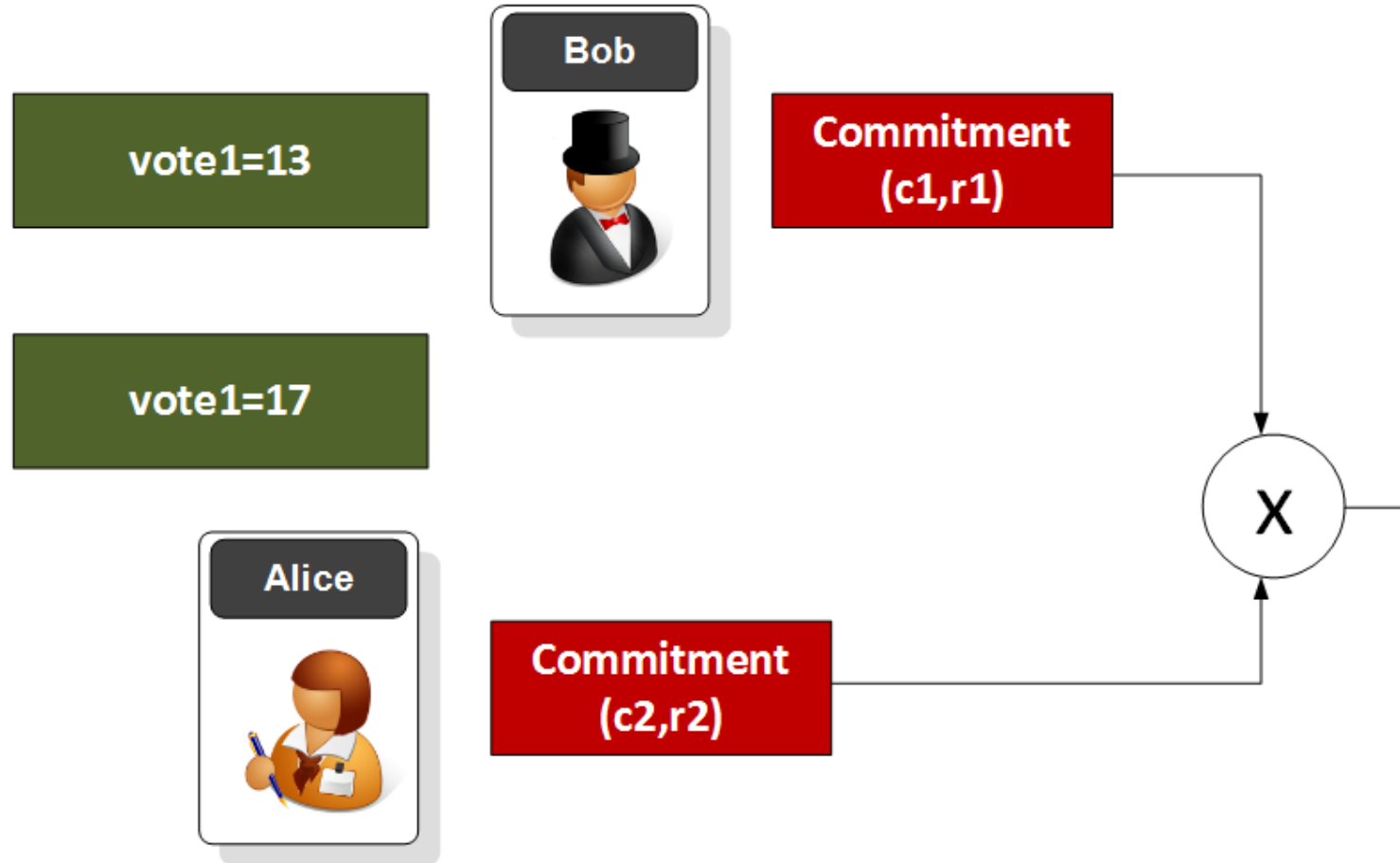
Pedersen Commitment: Voting



Pedersen Commitment: Homomorphic Encryption



Pedersen Commitment: Homomorphic Encryption



```
def add(self, param, *cm):  
    addCM = 1  
    for x in cm:  
        addCM *= x  
    addCM = addCM % param[1]  
    return addCM
```

```
class prover:  
    def commit(self, param, x):  
        q,g,h = generate(param)  
  
        r = number.getRandomRange(1, q-1)  
        c = (pow(g,x,q) * pow(h,r,q)) % q  
        return c, r
```

```
security = 80  
msg1 = 1  
msg2 = 2  
  
if (len(sys.argv)>1):  
    msg1=int(sys.argv[1])  
  
if (len(sys.argv)>2):  
    msg2=int(sys.argv[2])
```

```
v = verifier()  
p = prover()  
  
param = v.setup(security)  
  
c1, r1 = p.commit(param, msg1)  
c2, r2 = p.commit(param, msg2)  
  
addCM = v.add(param, c1, c2)
```

Bob



Alice



Pedersen Commitment

Prof Bill Buchanan, The Cyber Academy

<http://asecuritysite.com>

Eve



**CYBER
ACADEMY**